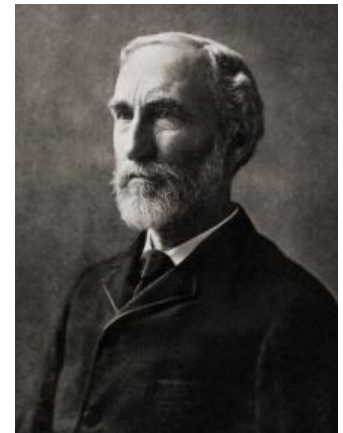


If you squint hard enough,

Gibbs distributions behave  
like mixtures of product  
measures



**Renan Gross, WIS**



**Ronen Eldan, WIS**

- Ok, so I need to tell you:
  - What Gibbs distributions are
  - What mixtures of product measures are
  - What it means to squint hard enough
  - What this is good for (maybe)

A probability measure  $X_n^f$  on  $\{-1,1\}^n$  is a

***Gibbs distribution with Hamiltonian  $f$***

if for  $X \sim X_n^f$ , we have

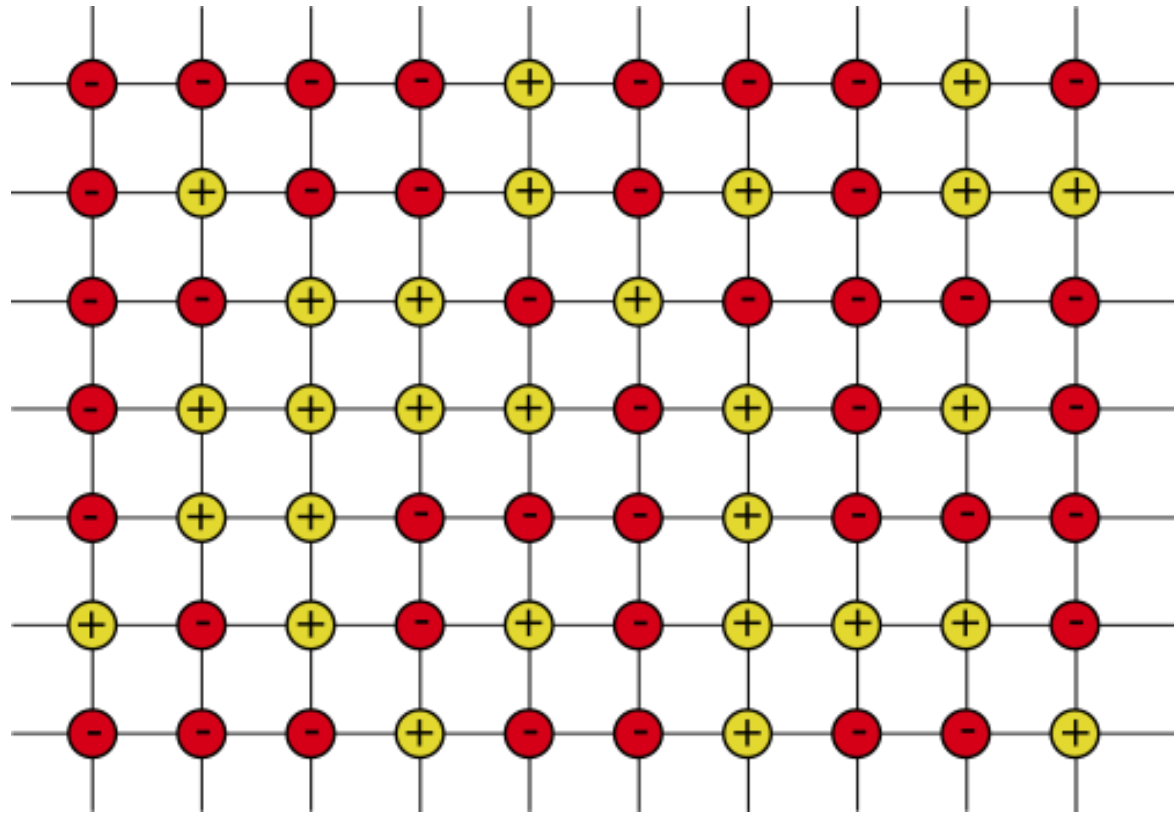
$$\mathbf{P}[X = \mathbf{x}] = \mathbf{e}^{f(\mathbf{x})} / \mathbf{Z}$$

Where  $Z$  is a normalizing constant.

Notes:

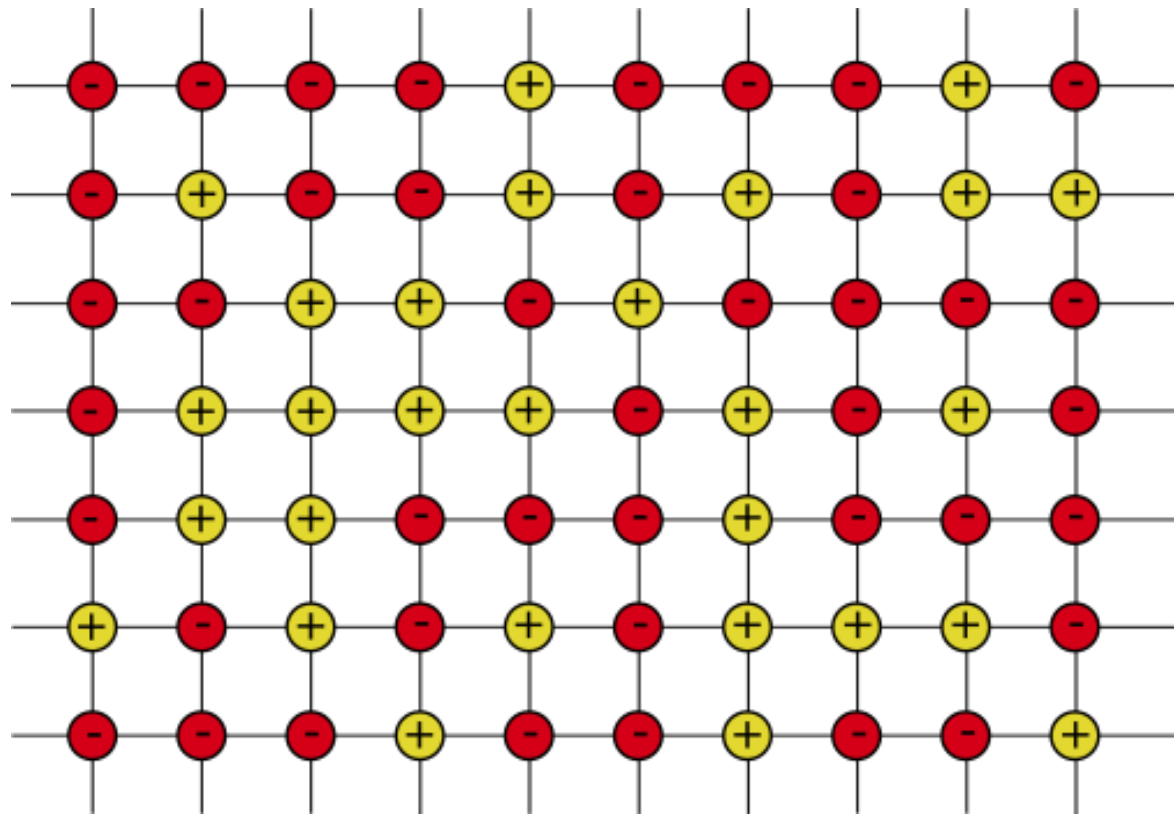
- $f: \{-1,1\}^n \rightarrow [-\infty, \infty]$
- Every distribution can be written this way...

# Ising Model



$$f(x) = -\beta J \sum_{i \sim j} x_i x_j + \beta H \sum_i x_i$$

# Ising Model



$$f(x) = \langle x, Ax \rangle + \langle \mu, x \rangle; \quad A \in \mathbb{R}^{n \times n}$$
$$\mu \in \mathbb{R}^n$$

# Exponential random graphs

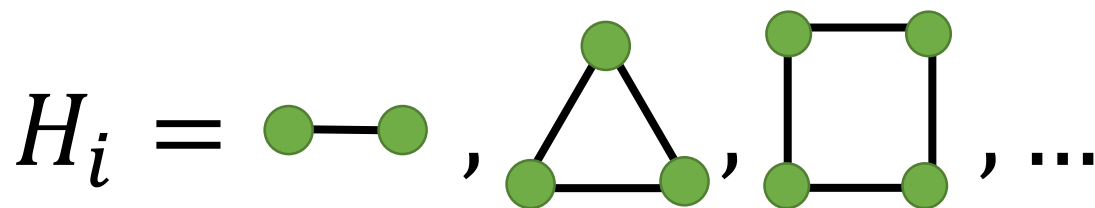
aka Graph Counting

Here,  $n = \binom{N}{2}$  represents edges over graphs with  $N$  vertices

Each **edge** is an entry in the vector

$$f(G) = \sum_i \beta_i \cdot \#\{\text{copies of } H_i \text{ in } G\}$$

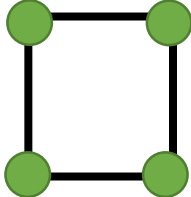
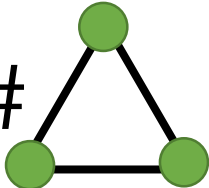
where



# Exponential random graphs

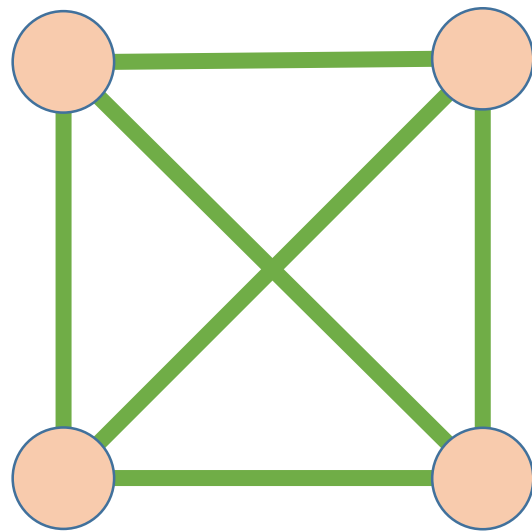
aka Graph Counting

For example, if

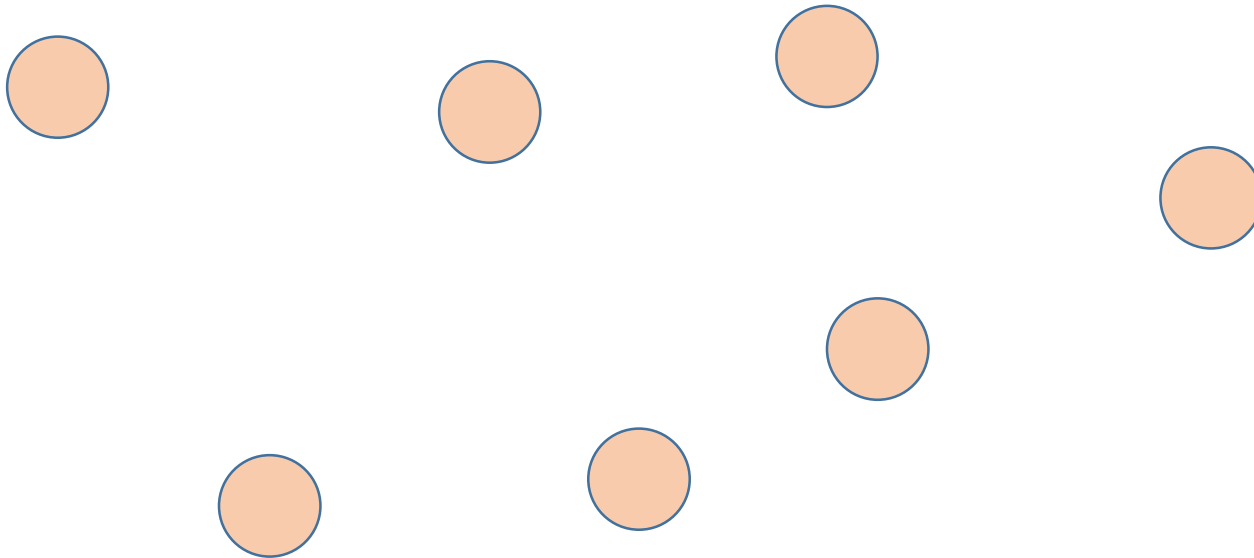
$$f(G) = -\beta_1 \# \text{triangle} + \beta_2 \# \text{square}$$


then

$$f(\text{square with diagonals}) = -4\beta_1 + 2\beta_2$$



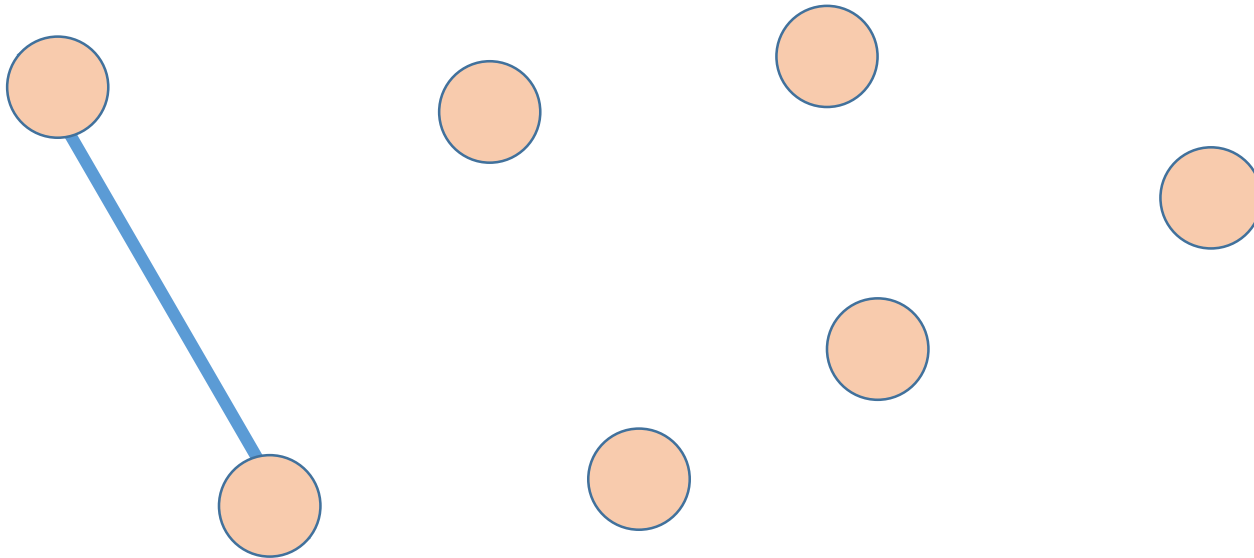
- Exponential random graphs are hard to work with.
- For example, when generating exponential random graphs with Glauber dynamics, many combinations of positive  $\beta$ 's give an exponential mixing time.



- For other  $\beta$ 's the mixing time is polynomial, but those correspond to uninteresting cases!

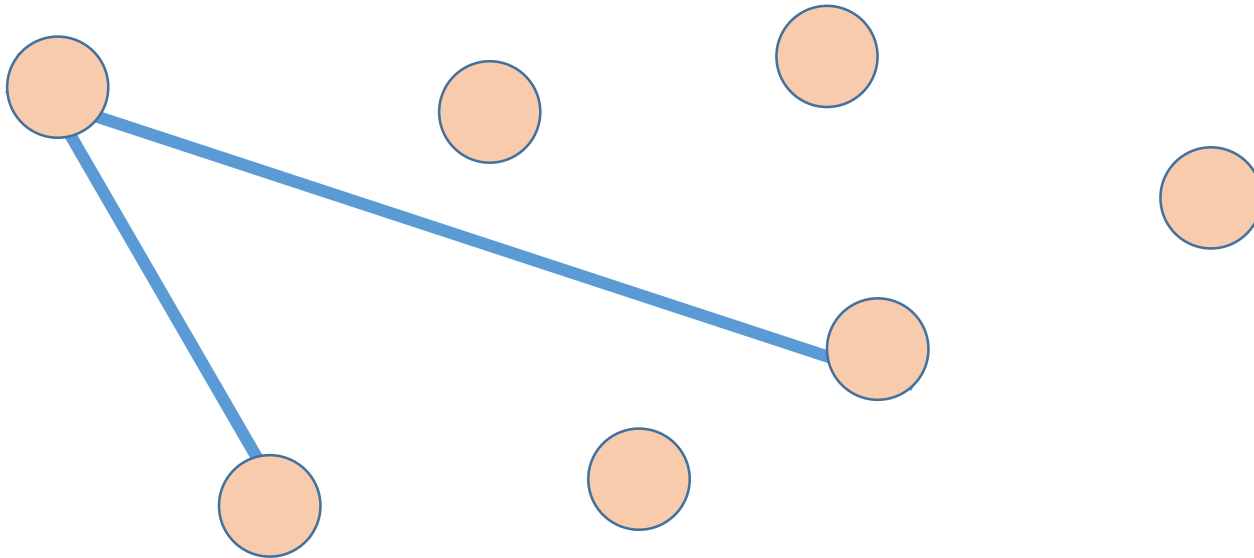


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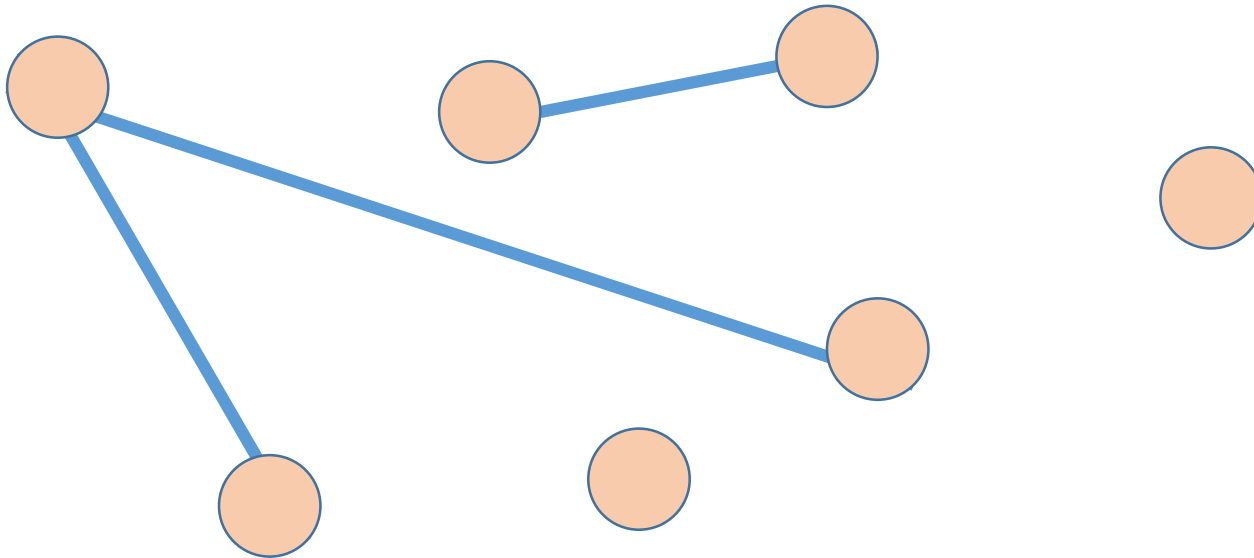
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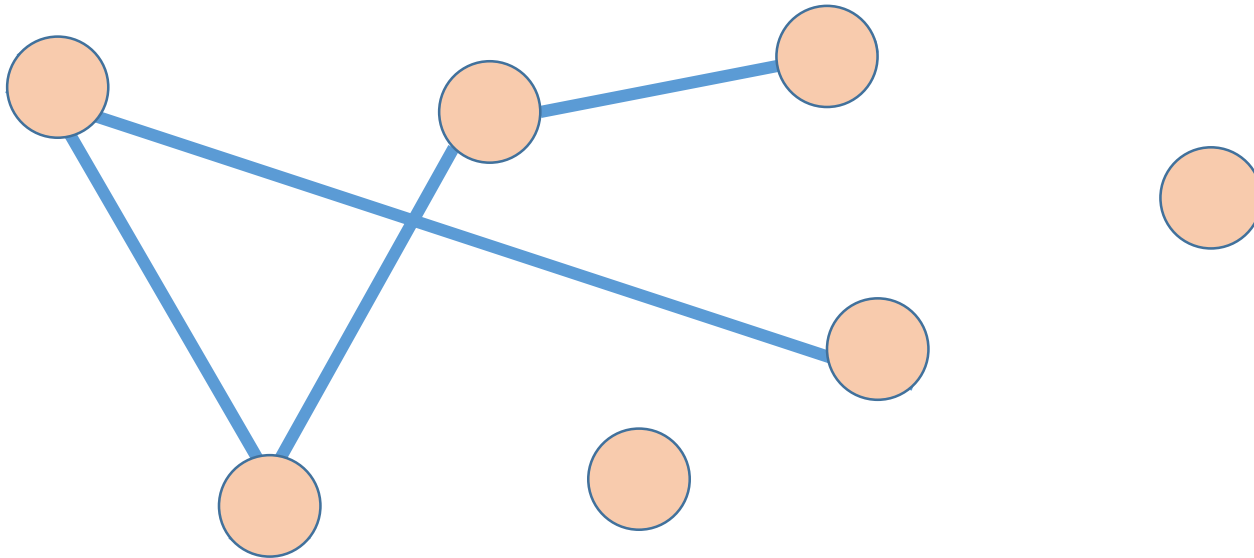
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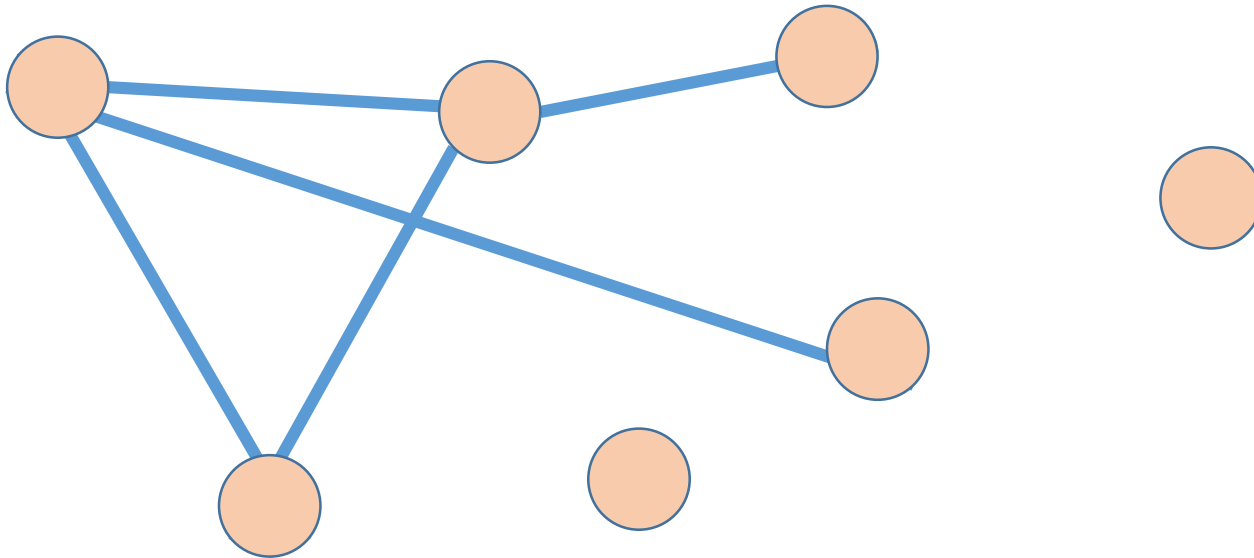
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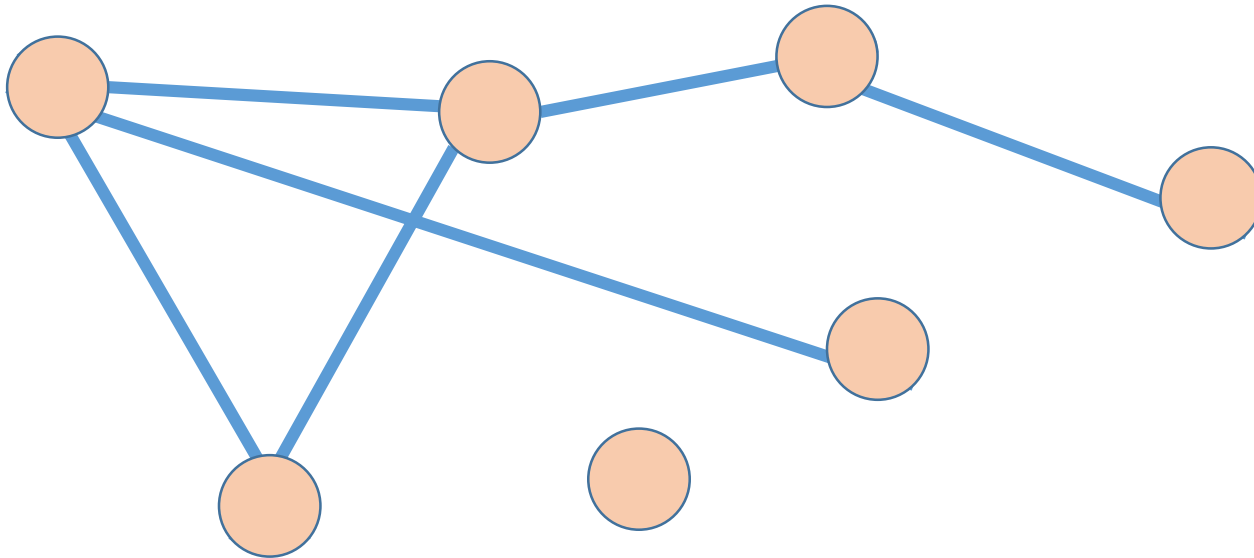
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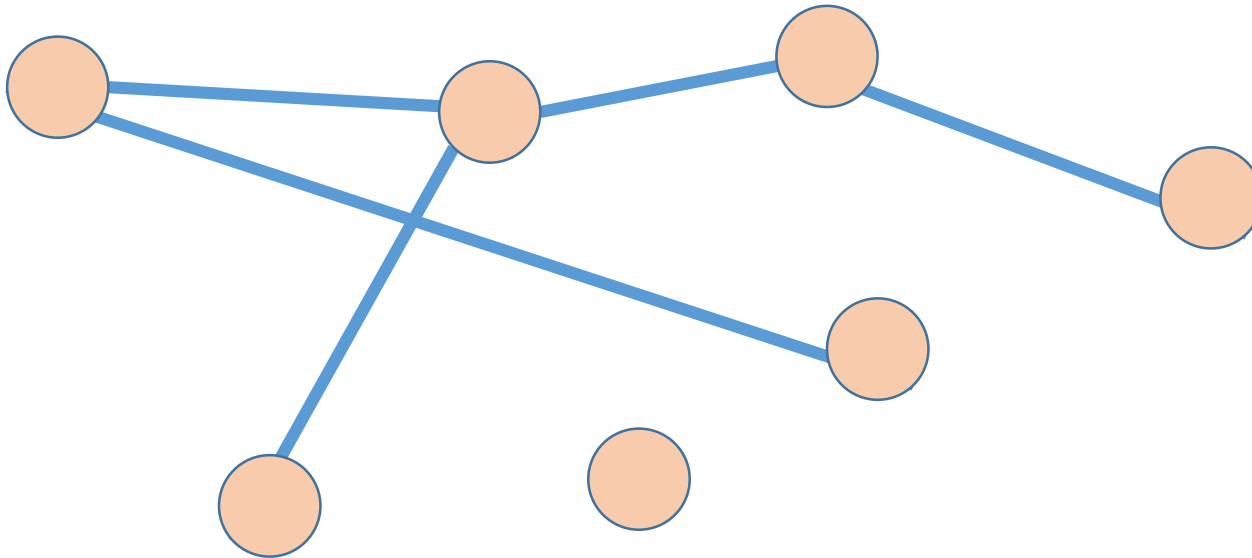
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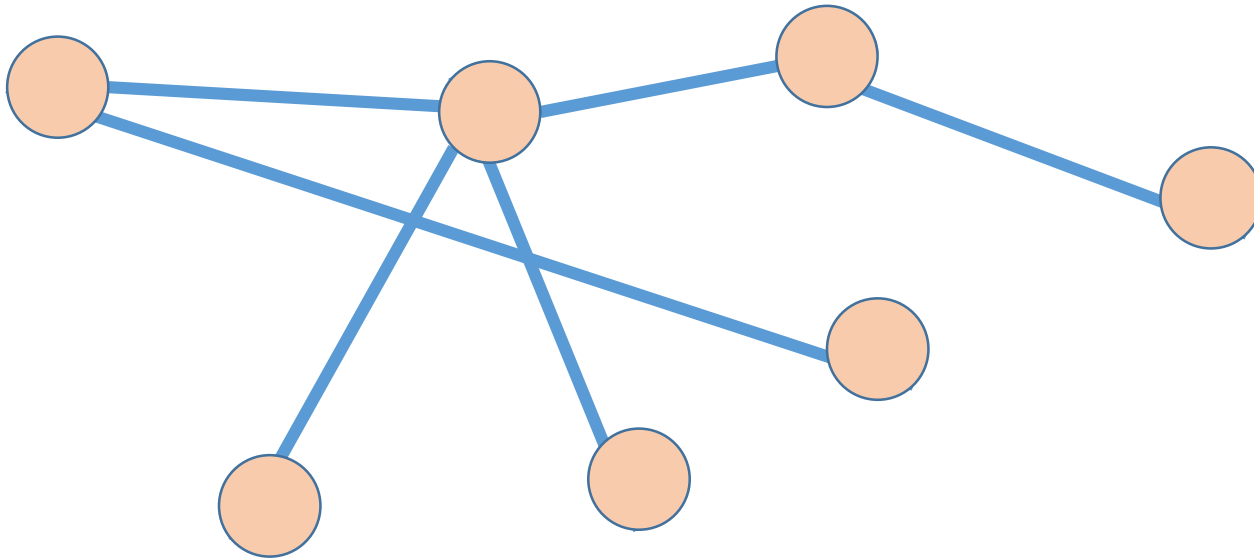
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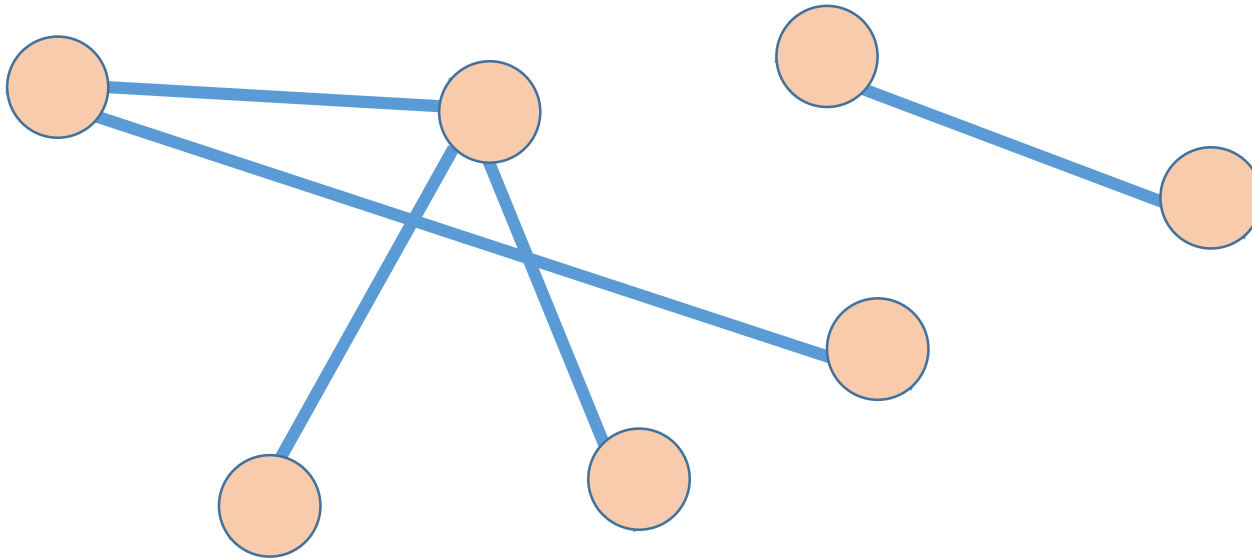
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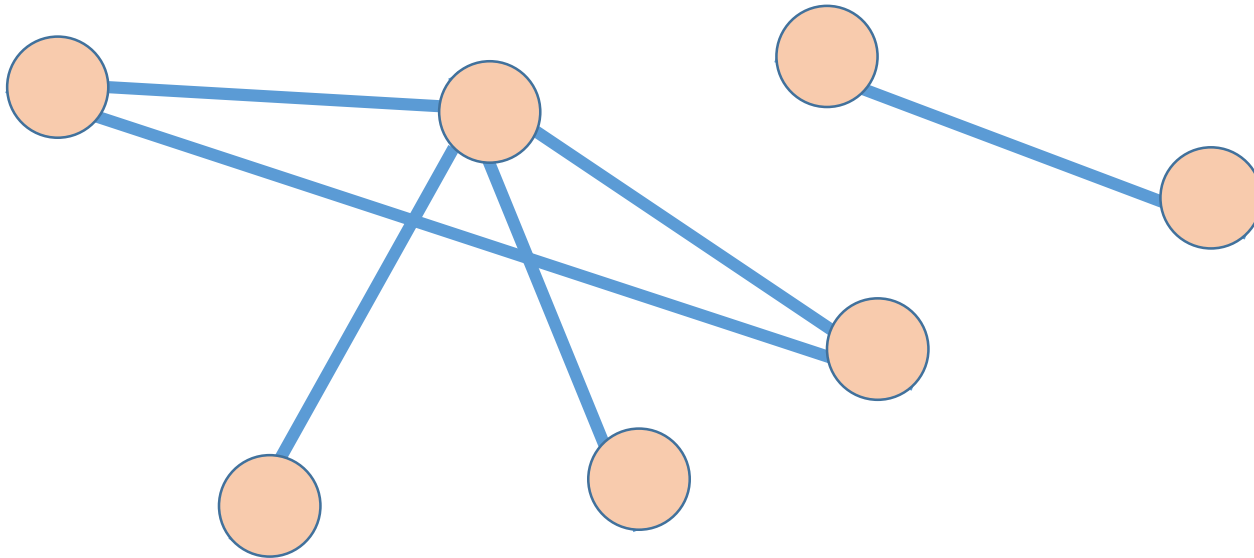


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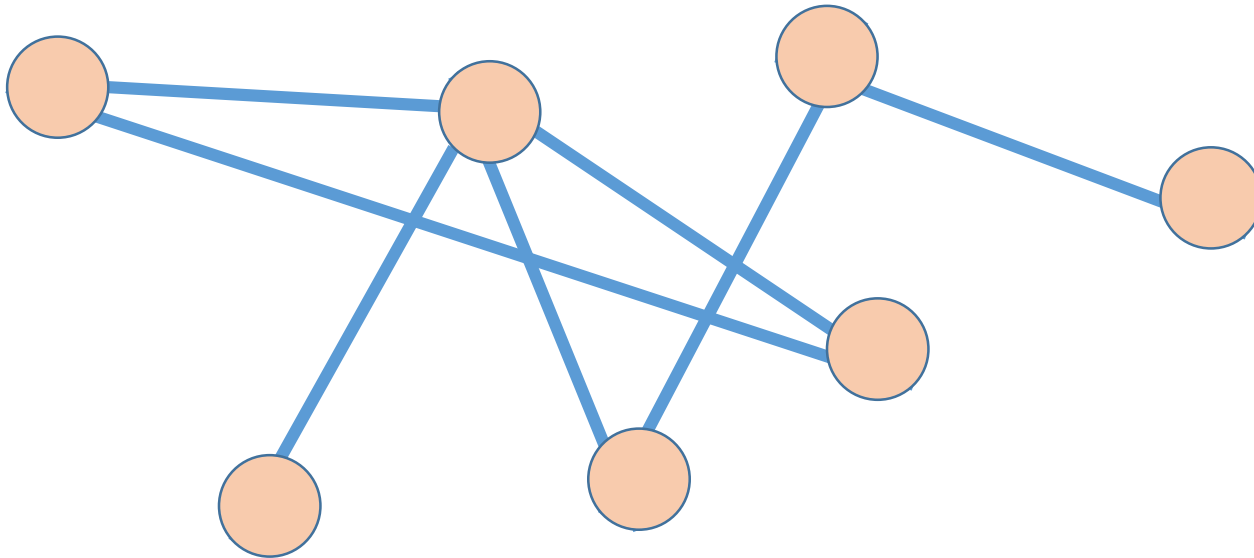
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A product measure is a measure where every entry is independent.

In a mixture of product measures, first we choose probabilities, then we choose from that probability.

Let  $\rho$  be a measure on  $[-1,1]^n$ . A random variable  $X(\rho)$  is a  $\rho$ -mixture if

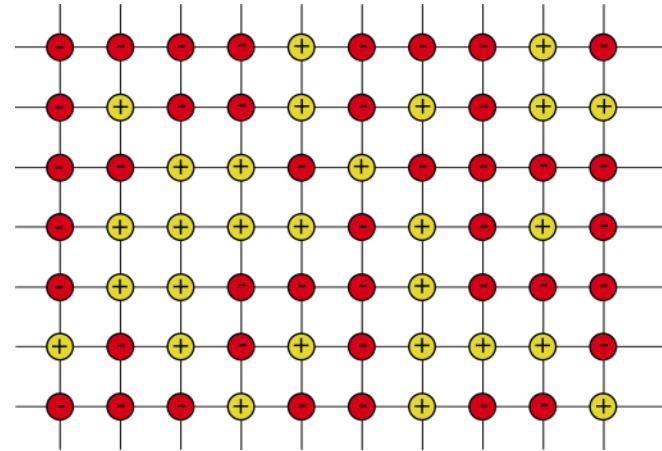
- $\Pr[X(\rho) = x] = \int \Pr[X(z) = x]d\rho(z)$

Where  $X(z)$  is the random vector with independent entries whose expectation is  $Z$ .

(Every distribution can be written as a mixture...)

For example, for Ising models:

$$f(x) = -\beta J \sum_{i \sim j} x_i x_j$$



We expect that:

- For small  $\beta$ , spins will be iid Bernoulli with probability  $\frac{1}{2}$
- For large  $\beta$ , spins tend to point either up or down.
  - We could expect that w.p  $\frac{1}{2}$ , spins are iid Bernoulli w.p  $p$ , and w.p  $\frac{1}{2}$ , spins are iid Bernoulli w.p  $(1-p)$ .

## Theorem:

$$P[X = x] = e^{f(x)} / Z$$

If the Hamiltonian  $f$  is “nice enough”, then

- The Gibbs distribution can be approximated by a  $\rho$ -mixture
- Further,  $\rho$  is concentrated on vectors  $X \in [-1,1]^n$  which satisfy

$$\|X - \tanh(\nabla f(X))\|_1 = o(n)$$

- Here,  $\partial_i f(y) = \frac{f(y^{i+}) - f(y^{i-})}{2}$ , i.e energy difference in flipping one coordinate for discrete vectors.
- Harmonic extension for continuous vectors.

- What is a nice Hamiltonian?
- “A nice Hamiltonian is a smooth Hamiltonian”



- What is a nice Hamiltonian?
- “A nice Hamiltonian is a smooth Hamiltonian”
- Borrowing from physics, you might call this a “mean-field” Hamiltonian
- A sufficient condition: the Gaussian width of the gradient of the Hamiltonian is low.

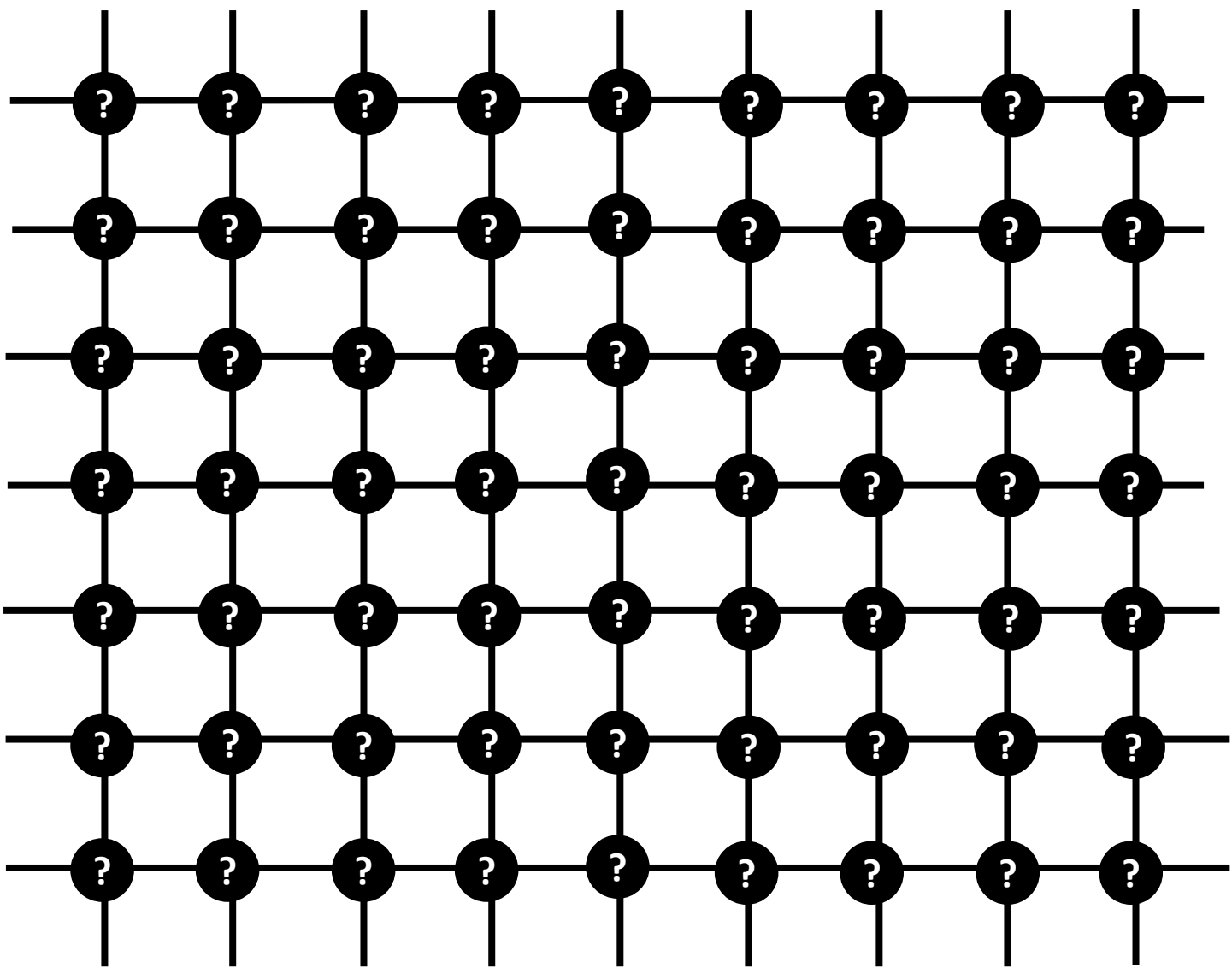


$$GW(K) = \mathbb{E} \left[ \sup_{x \in K} \langle x, \Gamma \rangle \right]$$

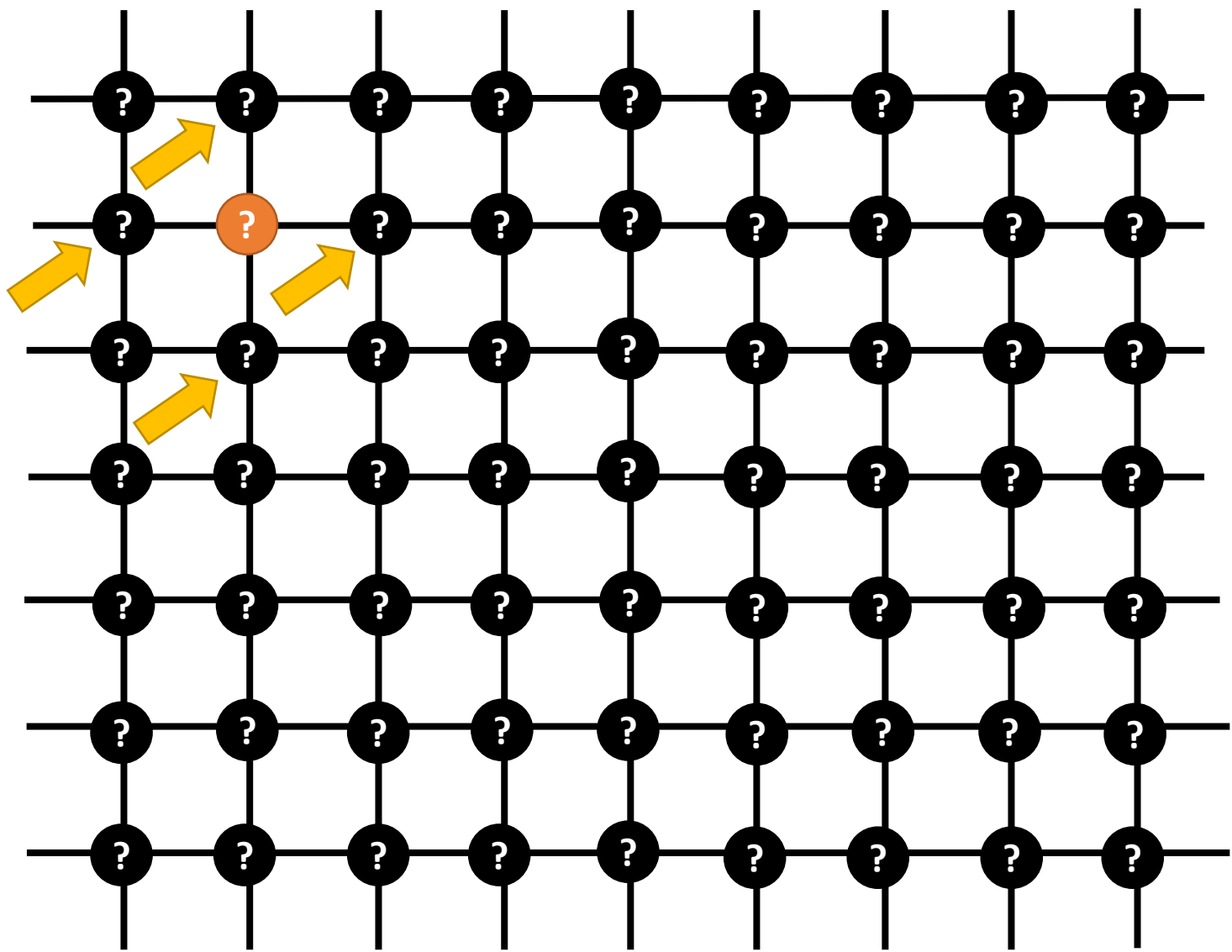
where  $\Gamma \sim N(0, \text{Id})$  is a standard Gaussian in  $\mathbb{R}^n$ .

- Correlation with random noise



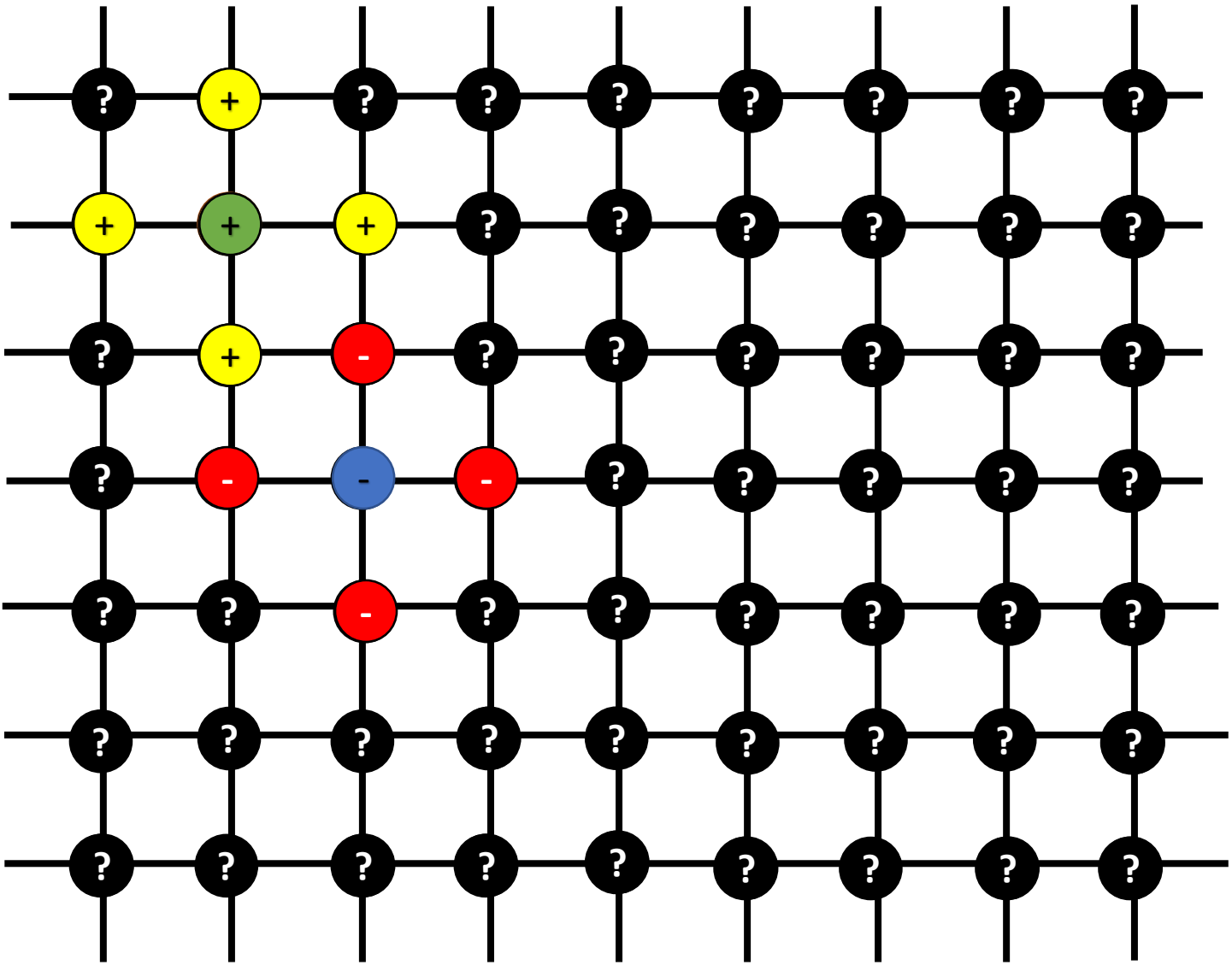








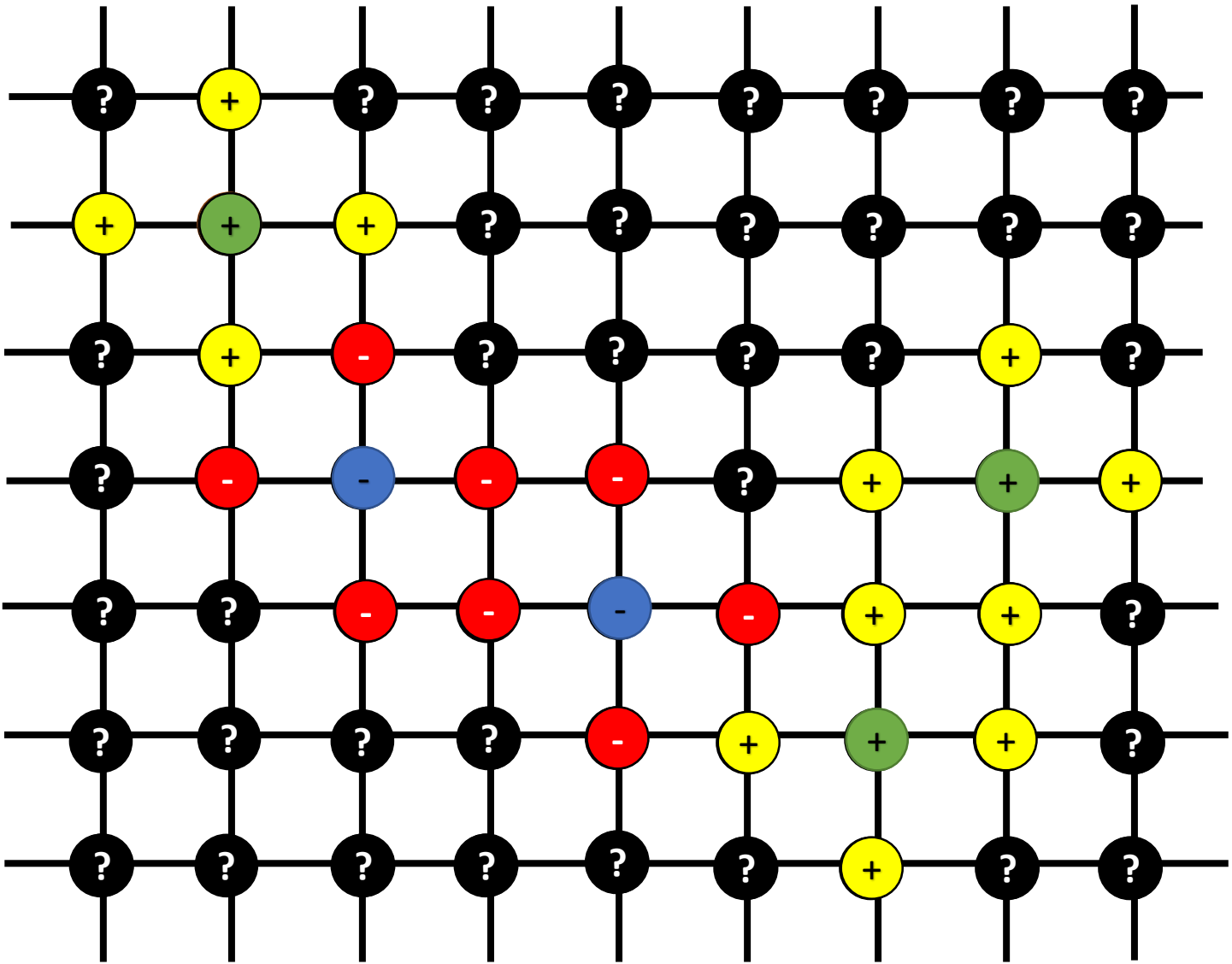


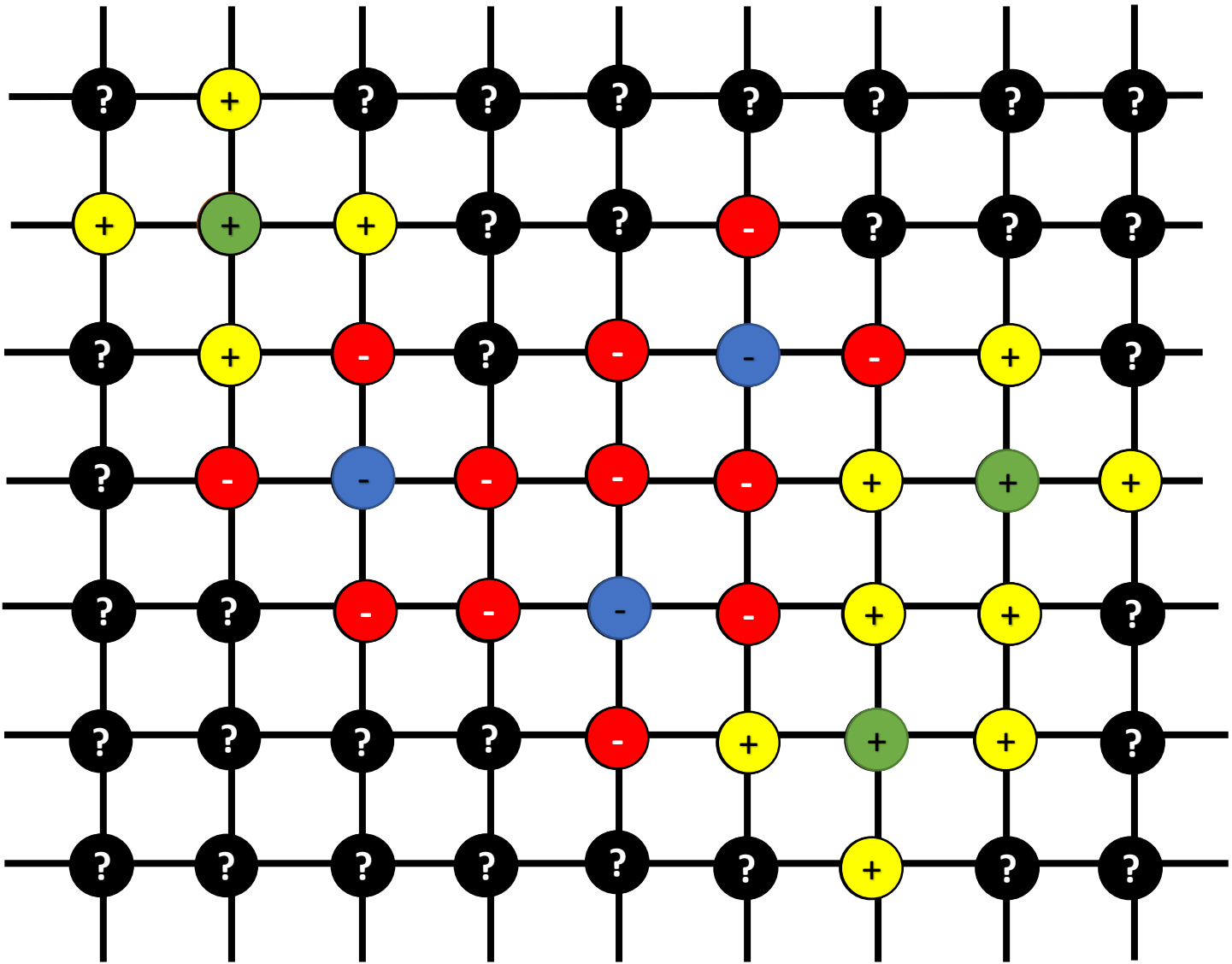


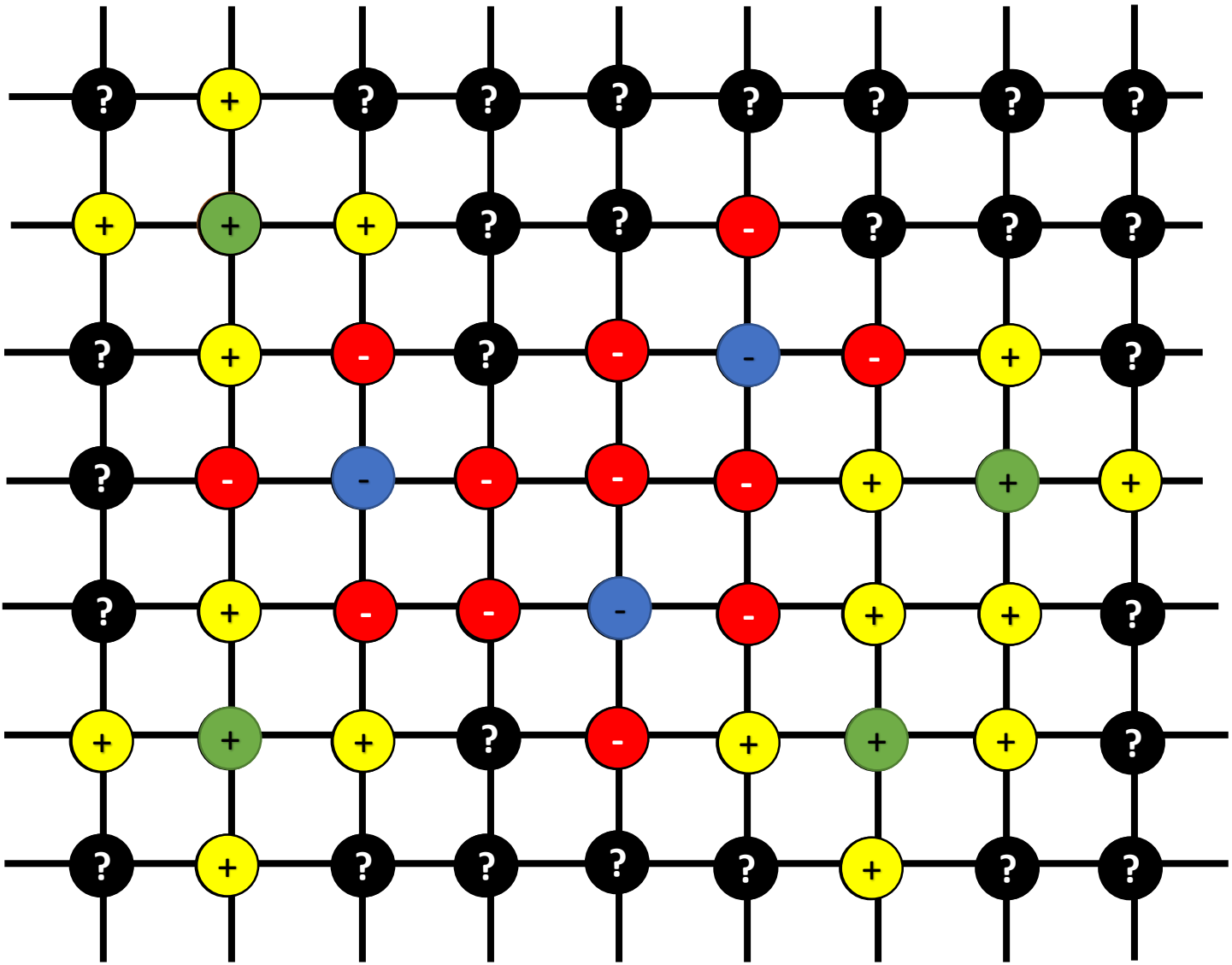


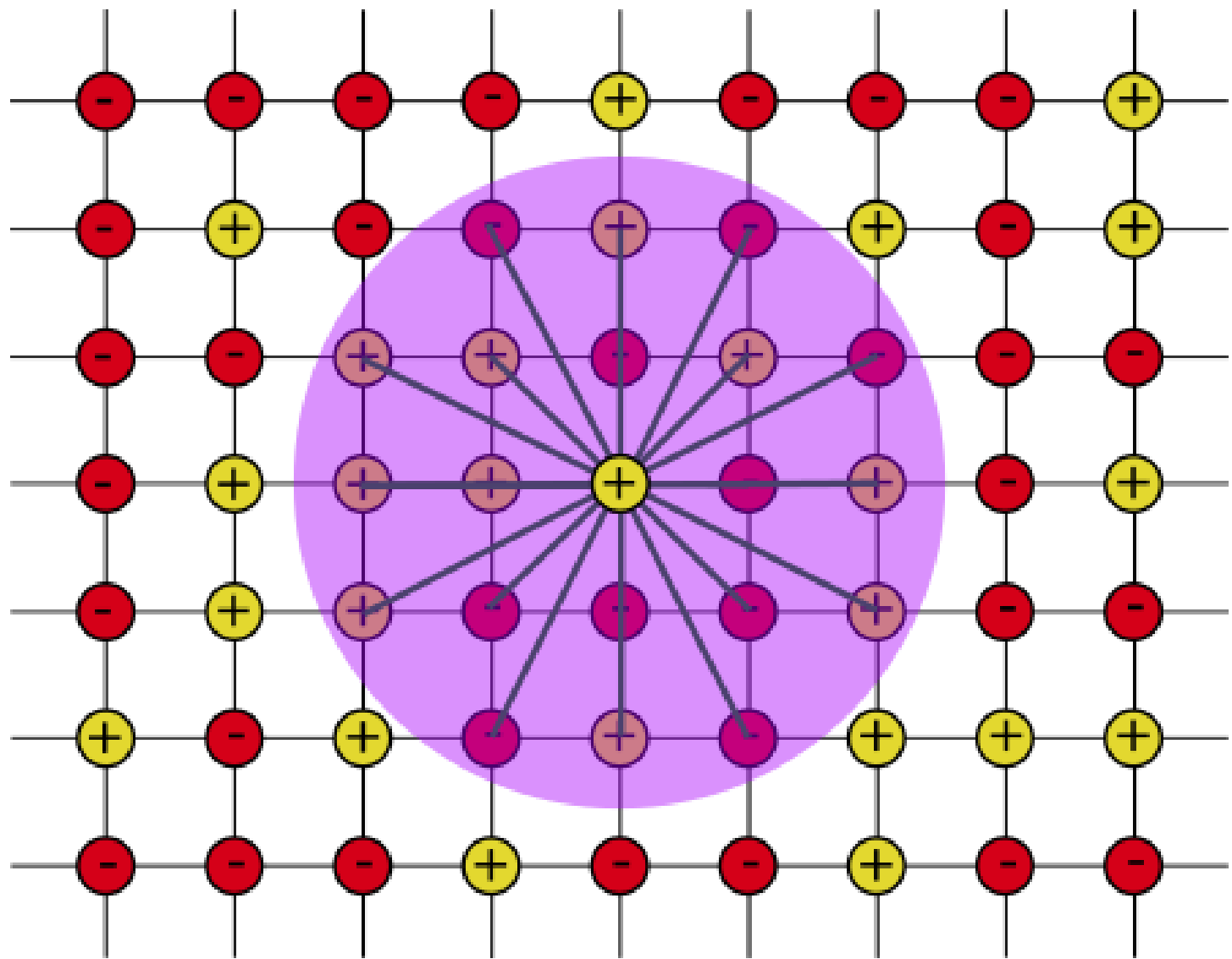


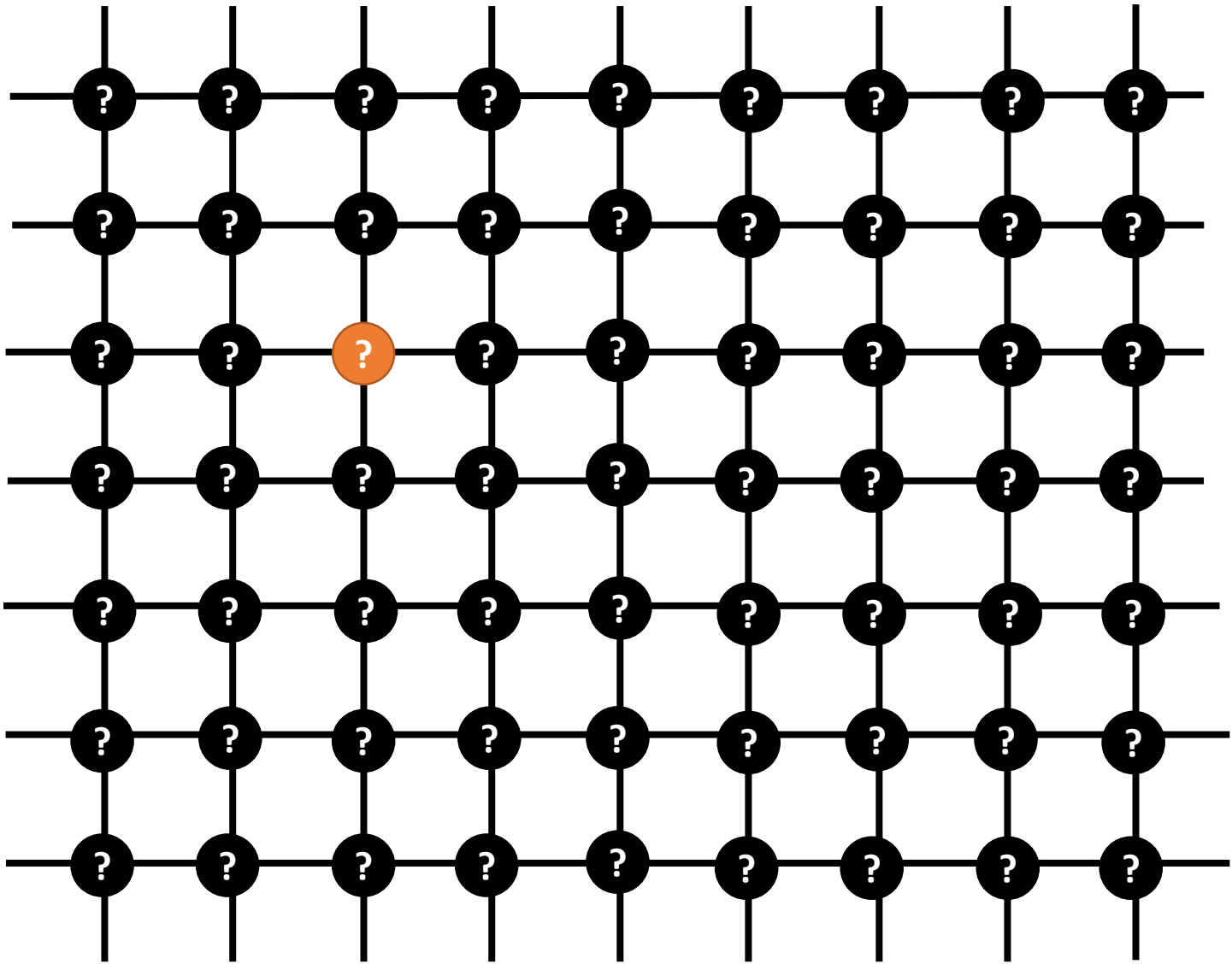
















- This equation

$$\|X - \tanh(\nabla f(X))\|_1 = o(n)$$

tells you how to break the symmetry of the system.

i.e. If there is a phase transition, you should find it in the solutions of the equation



Easy example:

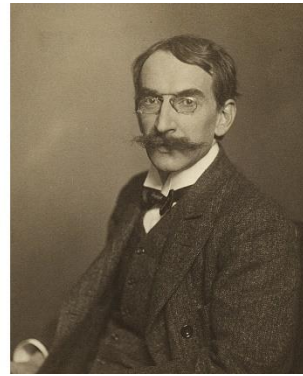
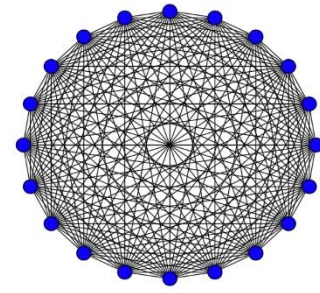
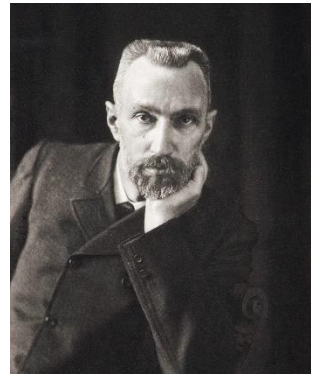
The Curie Weiss Ising model:

$$f(X) = \frac{\beta}{n} \sum_{i \neq j} x_i x_j$$

$$= x^T A x \text{ for } A = \begin{bmatrix} 0 & \dots & \frac{\beta}{n} \\ \vdots & \ddots & \vdots \\ \frac{\beta}{n} & \dots & 0 \end{bmatrix}$$

This gives  $\nabla f(x) = Ax$ , and so

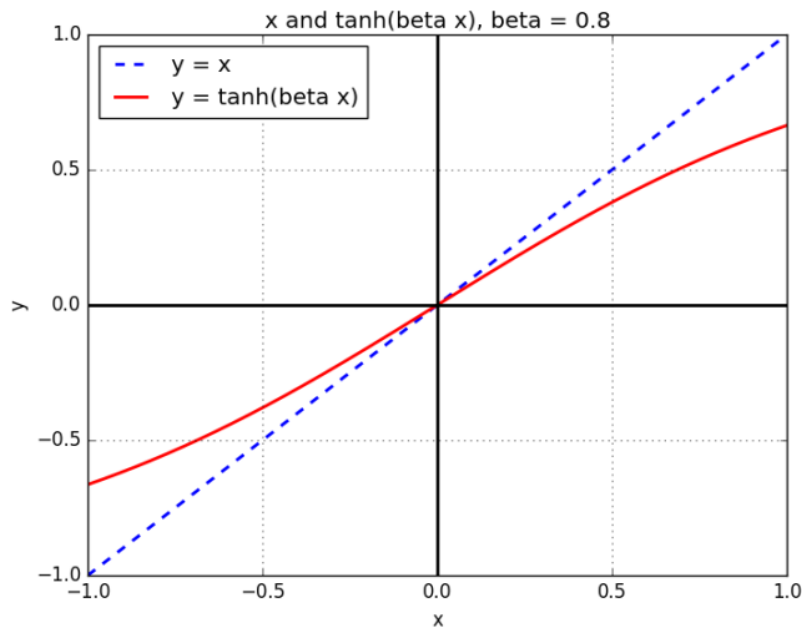
$$\|x - \tanh(Ax)\|_1 = o(n)$$



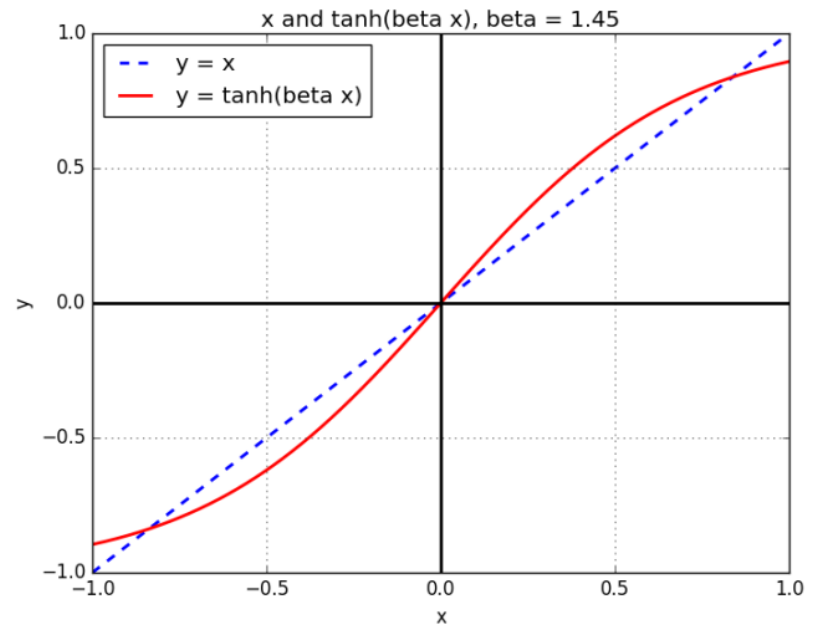
$$\|x - \tanh(Ax)\|_1 = o(n)$$

If  $x = \mathbf{1} \cdot \bar{x}$  is a scalar vector, we get a scalar equation

$$\bar{x} = \tanh(\beta \bar{x})$$



$$\beta < 1$$

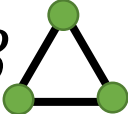


$$\beta > 1$$

- This equation

$$\|X - \tanh(\nabla f(X))\|_1 = o(n)$$

tells you how to break symmetry.

Harder example:  $-\beta$   counts in ERGM,

$$X = \frac{1 - \tanh\left(\frac{\beta}{n} X^2\right)}{2}$$

$$X = \frac{1 - \tanh\left(\frac{\beta}{n} X^2\right)}{2}$$

- For small  $\beta$ , only the trivial solution exists:

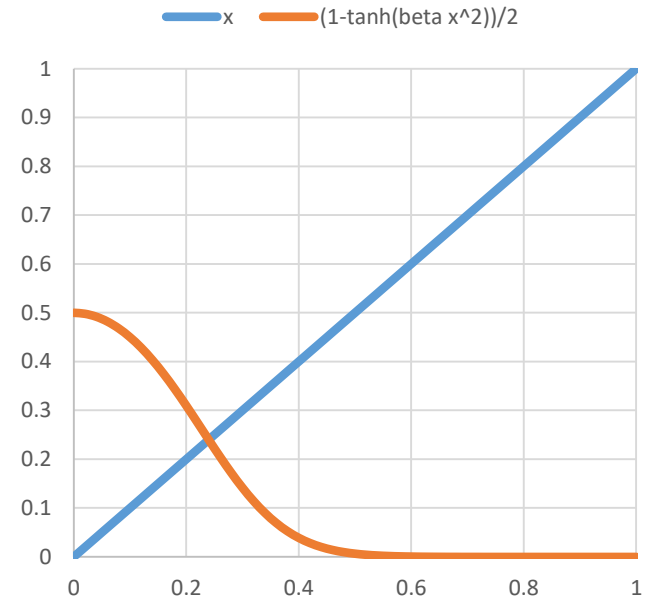
$$X = \mathbf{1} \cdot c$$

- For large  $\beta$  there exists a bipartite solution

- Corresponds to cuts in graphs!

- As  $\beta \rightarrow \infty$ , can tend to  $G\left(\frac{n}{2}, \frac{n}{2}, \frac{1}{2}\right)$

- What else?



For more information, call

1-800-<https://arxiv.org/abs/1708.05859>

Thanks!

